

# Independent partitions and indiscernibility

Akito Tsuboi

University of Tsukuba

Kyoto RIMS, 2010

November 29

# Outline

- 1 Simplicity and Independent Partitions,
  - 1 Definitions
  - 2 Examples
- 2 Ranks
  - 1  $D(\Sigma, \varphi, k)$
  - 2  $D(\Sigma, \varphi)$
  - 3  $D_{\text{inp}}$
- 3 Main Result

# Settings

- $T$  is a complete theory formulated in  $L$ .

# Settings

- $T$  is a complete theory formulated in  $L$ .
- We work in a very saturated  $\mathcal{M} \models T$ .

# Settings

- $T$  is a complete theory formulated in  $L$ .
- We work in a very saturated  $\mathcal{M} \models T$ .
- $a, b, \dots$  are (finite) tuples in  $\mathcal{M}$ .

# Settings

- $T$  is a complete theory formulated in  $L$ .
- We work in a very saturated  $\mathcal{M} \models T$ .
- $a, b, \dots$  are (finite) tuples in  $\mathcal{M}$ .
- $A, B, \dots$  are small sets in  $\mathcal{M}$ .

# Settings

- $T$  is a complete theory formulated in  $L$ .
- We work in a very saturated  $\mathcal{M} \models T$ .
- $a, b, \dots$  are (finite) tuples in  $\mathcal{M}$ .
- $A, B, \dots$  are small sets in  $\mathcal{M}$ .
- $I, J$  are sequences of tuples in  $\mathcal{M}$ .

# Settings

- $T$  is a complete theory formulated in  $L$ .
- We work in a very saturated  $\mathcal{M} \models T$ .
- $a, b, \dots$  are (finite) tuples in  $\mathcal{M}$ .
- $A, B, \dots$  are small sets in  $\mathcal{M}$ .
- $I, J$  are sequences of tuples in  $\mathcal{M}$ .
- $M, N, \dots \prec \mathcal{M}$ .



# Settings

- $T$  is a complete theory formulated in  $L$ .
- We work in a very saturated  $\mathcal{M} \models T$ .
- $a, b, \dots$  are (finite) tuples in  $\mathcal{M}$ .
- $A, B, \dots$  are small sets in  $\mathcal{M}$ .
- $I, J$  are sequences of tuples in  $\mathcal{M}$ .
- $M, N, \dots \prec \mathcal{M}$ .
- Formulas are denoted by  $\varphi, \psi, \dots$

# Settings

- $T$  is a complete theory formulated in  $L$ .
- We work in a very saturated  $\mathcal{M} \models T$ .
- $a, b, \dots$  are (finite) tuples in  $\mathcal{M}$ .
- $A, B, \dots$  are small sets in  $\mathcal{M}$ .
- $I, J$  are sequences of tuples in  $\mathcal{M}$ .
- $M, N, \dots \prec \mathcal{M}$ .
- Formulas are denoted by  $\varphi, \psi, \dots$
- $m, n, k, \dots$  are natural numbers.

# Simple Theory

A **simple** theory is characterized as a theory in which the length of dividing sequence of types is bounded ( $< \infty$ ).

# Low Theory

A **low** theory is characterized by the following property: For each formula  $\varphi(x, y)$  there is a number  $n_\varphi \in \omega$  such that whenever  $\{\varphi(x, a_i) : i < m\}$  satisfies

1  $\{\varphi(x, a_i) : i < m\}$  is consistent, and

2  $\varphi(x, a_i)$  divides over  $A_i = \{a_j : j < i\}$  ( $i < m$ ),

then  $m \leq n_\varphi$ .

# Non-Low Simple Theory

Casanovas constructed a simple nonlow theory  
 $T_1 = Th(M, P, P_1, P_2, \dots, Q, R)$ .

- 1  $M$  is the disjoint union of  $P$  and  $Q$ .

- 1  $M$  is the disjoint union of  $P$  and  $Q$ .
- 2  $P_1, P_2, \dots$  are disjoint copies of  $\omega$ .

- 1  $M$  is the disjoint union of  $P$  and  $Q$ .
- 2  $P_1, P_2, \dots$  are disjoint copies of  $\omega$ .
- 3  $P = \bigcup_{i \in \omega} P_i \cup G$ , where  $G$  is a random graph.



- 1  $M$  is the disjoint union of  $P$  and  $Q$ .
- 2  $P_1, P_2, \dots$  are disjoint copies of  $\omega$ .
- 3  $P = \bigcup_{i \in \omega} P_i \cup G$ , where  $G$  is a random graph.
- 4  $Q$  is the set of all sequences  $(A_1, A_2, \dots, A_\omega)$ , where  $A_n$  is an  $n$ -element subset of  $P_n$  and for some  $a \in G$ ,  $A_\omega \subset G$  is the set of all  $g \in G$  directly connected to  $a$ .

- 1  $M$  is the disjoint union of  $P$  and  $Q$ .
- 2  $P_1, P_2, \dots$  are disjoint copies of  $\omega$ .
- 3  $P = \bigcup_{i \in \omega} P_i \cup G$ , where  $G$  is a random graph.
- 4  $Q$  is the set of all sequences  $(A_1, A_2, \dots, A_\omega)$ , where  $A_n$  is an  $n$ -element subset of  $P_n$  and for some  $a \in G$ ,  $A_\omega \subset G$  is the set of all  $g \in G$  directly connected to  $a$ .
- 5  $R \subset P \times Q$ .
- 6  $R(a, (A_1, A_2, \dots, A_\omega))$  if (i)  $a \in P_n$  and  $a \in A_n$  ( $\exists n \in \omega$ ) or (ii)  $a \notin \bigcup_n P_n$  and  $a \in A_\omega$ .

This theory  $T_1$  is not supersimple.  $R(x, y)$  defines infinitely many mutually **independent partitions** in the following sense: If we enumerate  $P_n$  as  $P_n = \{a_{nm} : m \in \omega\}$ , then

This theory  $T_1$  is not supersimple.  $R(x, y)$  defines infinitely many mutually **independent partitions** in the following sense: If we enumerate  $P_n$  as

$P_n = \{a_{nm} : m \in \omega\}$ , then

- for each  $\eta \in \omega^\omega$ ,  $\{R(a_{n\eta(n)}, y) : n = 1, 2, \dots\}$  is consistent, and
- for each  $n = 1, 2, \dots$ ,  $\{R(a_{nm}, y) : m \in \omega\}$  is  $(n + 1)$ -inconsistent.

# Non-Low Supersimple Theory

By modifying  $T_1$ , Casanovas and Kim showed the existence of a supersimple nonlow theory  $T_2$ . This  $T_2$  does not have infinitely many mutually independent partitions.

# Non-Low Supersimple Theory

By modifying  $T_1$ , Casanovas and Kim showed the existence of a supersimple nonlow theory  $T_2$ . This  $T_2$  does not have infinitely many mutually independent partitions.

However, for each  $k \in \omega$ , we can find a formula  $\varphi(x, y)$  and parameter sets  $A_i = \{a_{ij} : j \in \omega\}$  ( $i < k$ ) defining  $k$  independent partitions.

$D_{\text{inp}}(*, *)$ 

## Definition

$D_{\text{inp}}(\Sigma(x), \varphi(x, y))$  is the first cardinal  $\kappa$  such that there are no  $\kappa$ -many independent partitions  $\{\varphi(x, a_{ij}) : j \in \omega\}$  ( $i < \kappa$ ) of  $\Sigma$ .

## Remark

- For  $T_1$ ,  $D_{\text{inp}}(x = x, R(y, x)) = \omega_1$ .
- For  $T_2$ , for some  $\varphi(x, y)$ ,  
 $D_{\text{inp}}(x = x, \varphi(x, y)) = \omega$ .



So it is natural to ask whether there is a simple nonlow theory  $T$  such that

$$D_{\text{inp}}(x = x, \varphi(x, y)) < \omega,$$

for any  $\varphi$ .

# Ranks

First we recall definitions of basic ranks.

Let  $\Sigma(x)$  be a set of formulas and  $\varphi(x, y)$  a formula.

Let  $k \in \omega$ .

## Definition

$D(\Sigma(x), \varphi(x, y), k)$

- 1
  - $D(\Sigma(x), \varphi(x, y), k) \geq 0$  if  $\Sigma(x)$  is consistent.
  - $D(\Sigma(x), \varphi(x, y), k) \geq n + 1$  if there is an indiscernible sequence  $\{b_i : i \in \omega\}$  over  $\text{dom}(\Sigma)$  such that  $D(\Sigma(x) \cup \{\varphi(x, b_i)\}, \varphi(x, y), k) \geq n$  for all  $i \in \omega$ , and  $\{\varphi(x, b_i) : i \in \omega\}$  is  $k$ -inconsistent.

## Definition

- 2
- $D(\Sigma(x), \varphi(x, y)) \geq \mathbf{0}$  if  $\Sigma(x)$  is consistent.
  - For a limit ordinal  $\delta$ ,  $D(\Sigma(x), \varphi(x, y)) \geq \delta$  if  $D(\Sigma(x), \varphi(x, y)) \geq \alpha$  for all  $\alpha < \delta$ .
  - $D(\Sigma(x), \varphi(x, y)) \geq \alpha + \mathbf{1}$  if there is an indiscernible sequence  $\{b_i : i \in \omega\}$  over  $\mathbf{dom}(\Sigma)$  such that  $D(\Sigma(x) \cup \{\varphi(x, b_i)\}, \varphi(x, y)) \geq \alpha$  ( $i \in \omega$ ), and  $\{\varphi(x, b_i) : i \in \omega\}$  is inconsistent.

## Fact

- 1  $D(\Sigma(x), \varphi(x, y), k) \geq n$  if there is a tree  $A = \{a_v : v \in \omega^{<n}\}$  such that (1)  $\Sigma(x) \cup \{\varphi(x, a_{\eta|i}) : 1 \leq i \leq n\}$  is consistent ( $\forall \eta \in \omega^n$ ), and (2)  $\{\varphi(x, a_{v\sim i}) : i \in \omega\}$  is  $k$ -inconsistent ( $\forall v \in \omega^{<n}$ ).

## Fact

- 2  $D(\Sigma(x), \varphi(x, y)) \geq n$  if there is a tree  $A = \{a_v : v \in \omega^{<n}\}$  and numbers  $k_0, \dots, k_{n-1}$  such that (1)  $\Sigma(x) \cup \{\varphi(x, a_{\eta i}) : 1 \leq i \leq n\}$  is consistent ( $\forall \eta \in \omega^n$ ), and (2)  $\{\varphi(x, a_{v \smallfrown i}) : i \in \omega\}$  is  $k_{\text{lh}(v)}$ -inconsistent ( $\forall v \in \omega^{<n}$ ).

# Main Result

## Theorem

*Suppose that the size of independent partitions is bounded in  $\mathcal{T}$ . Then the following are equivalent:*

- 1  $\mathcal{T}$  is simple.*
- 2  $\mathcal{T}$  is low.*

## Proposition

*Suppose  $D_{\text{inp}}(x = x, \varphi(x, y)) = k - 1 < \omega$  and  $D(x = x, \varphi(x, y)) \geq \omega$ . Then  $T$  is not simple.*



# Proof.

Fix  $m \in \omega$ .

- By  $D(x = x, \varphi(x, y)) \geq \omega$ , there is a set  $A = \{a_\nu : \nu \in \omega^{\leq m}\}$  witnessing  $D(x = x, \varphi(x, y)) \geq m$ .

# Proof.

Fix  $m \in \omega$ .

- By  $D(x = x, \varphi(x, y)) \geq \omega$ , there is a set  $A = \{a_\nu : \nu \in \omega^{\leq m}\}$  witnessing  $D(x = x, \varphi(x, y)) \geq m$ .
- We have
  - 1  $\{\varphi(x, a_{\eta|i}) : 1 \leq i \leq m\}$  is consistent ( $\forall \eta \in \omega^m$ ),
  - 2  $\{\varphi(x, a_{\nu \frown i}) : i \in \omega\}$  is  $k_{\text{lh}(\nu)}$ -inconsistent ( $\forall \nu \in \omega^{< m}$ ).

- We can assume that  $A$  is an indiscernible tree.

- We can assume that  $A$  is an indiscernible tree.

- For  $\nu \in \omega^m$ , let  $\nu^*$  be the sequence

$$\nu(0), 0^k, \nu(1), 0^k, \dots, \nu(\text{lh}(\nu) - 1), 0^k.$$

- For  $\nu = \nu_0 \widehat{m}$ , let

$$a_\nu^* = a_{\nu_0^* \widehat{m} \widehat{0}}, a_{\nu_0^* \widehat{m} \widehat{0}^2}, \dots, a_{\nu_0^* \widehat{m} \widehat{0}^k}.$$

- Let  $\varphi^*(x, y_1, \dots, y_k)$  be the formula  $\varphi(x, y_1) \wedge \dots \wedge \varphi(x, y_k)$ .
- **Claim A**  $\{\varphi^*(x, a^*_{v_0 \widehat{m}}) : m \in \omega\}$  is  $k$ -inconsistent.

- Let  $\varphi^*(x, y_1, \dots, y_k)$  be the formula  $\varphi(x, y_1) \wedge \dots \wedge \varphi(x, y_k)$ .
- **Claim A**  $\{\varphi^*(x, a^*_{v_0 \widehat{\ } m}) : m \in \omega\}$  is  $k$ -inconsistent.

Suppose this is not the case. Then there is  $F = \{i_1, \dots, i_k\} \subset \omega$  such that

$$\{\varphi^*(x, a^*_{v_0 \widehat{\ } i_1}), \dots, \varphi^*(x, a^*_{v_0 \widehat{\ } i_k})\}$$

is consistent.

- By the definition of  $\varphi^*$ , in particular, the following set is consistent.

$$\{\varphi(x, a_{v_0^* \smallfrown i_1 \smallfrown 0}), \dots, \varphi(x, a_{v_0^* \smallfrown i_k \smallfrown 0^k})\}$$



- By the definition of  $\varphi^*$ , in particular, the following set is consistent.

$$\{\varphi(x, a_{v_0^* \hat{\sim} i_1 \hat{\sim} 0}), \dots, \varphi(x, a_{v_0^* \hat{\sim} i_k \hat{\sim} 0^k})\}$$

- For each  $\nu$  of length  $k$ , let  $\Gamma_\nu$  be the set:

$$\{\varphi(x, a_{v_0^* \hat{\sim} i_1 \hat{\sim} \nu(1)}), \dots, \varphi(x, a_{v_0^* \hat{\sim} i_k \hat{\sim} 0^{k-1} \hat{\sim} \nu(k)})\}.$$

- Then each  $\Gamma_\nu$  is consistent, by the indiscernibility of  $A$ .

- Then each  $\Gamma_\nu$  is consistent, by the indiscernibility of  $A$ .
- On the other hand, by our choice of the tree  $A$ , for each  $l = 0, \dots, k - 1$ , the set

$$\{\varphi(x, a^*_{\nu_0 \frown i_2 \frown 0^l \frown i}) : i \in \omega\}$$

is inconsistent ( $k_{\text{lh}(\nu_0) + (1+l)}$ -inconsistent).

- Then each  $\Gamma_\nu$  is consistent, by the indiscernibility of  $A$ .
- On the other hand, by our choice of the tree  $A$ , for each  $l = 0, \dots, k - 1$ , the set

$$\{\varphi(x, a^*_{\nu_0 \frown i_2 \frown 0^l \frown i}) : i \in \omega\}$$

is inconsistent ( $k_{\text{lh}(\nu_0) + (1+l)}$ -inconsistent).

- This yields  $D_{\text{inp}}(x = x, \varphi(x, z)) \geq k$ , a contradiction. (End of Proof of Claim)

- By Claim A, the set  $\{\varphi^*(x, a^*_\nu) : \nu \in \omega^m\}$  witnesses  $D(x = x, \varphi^*, k) \geq m$ .

- By Claim A, the set  $\{\varphi^*(x, a^*_\nu) : \nu \in \omega^m\}$  witnesses  $D(x = x, \varphi^*, k) \geq m$ .
- Since  $m$  is arbitrary, we conclude  $D(x = x, \varphi^*, k) = \infty$ , which means that  $T$  is not simple.