

proof $x \in I$, $y = f(x) \in J$ とする. 定理 1.2 の).
 $\exists \varphi_1(t), \exists \varphi_2(t) \wedge \exists \delta > 0$ s.t. $|h| < \delta$ の範囲で

$$f(x+h) = f(x) + f'(x)h + \varphi_1(h), \dots \textcircled{1}$$

$$g(y+h) = g(y) + g'(y)h + \varphi_2(h), \dots \textcircled{2}$$

$$\varphi_1(h) = o(h), \quad \varphi_2(h) = o(h).$$

→ δ . Δx に対して δ を選ぶ. $\textcircled{1}$ より

$$f(x+\Delta x) = f(x) + \frac{f'(x)(\Delta x) + \varphi_1(\Delta x)}{\Delta y}$$

\parallel
 Δy

と選ぶ.

$\Delta x \rightarrow 0$ のとき $\Delta y \rightarrow 0$ となる. $|\Delta x| \in \textcircled{1}$ の範囲で
 $|\Delta y| < \delta$ となるように Δx を選ぶ.

よって

$$\begin{aligned} (g \circ f)(x + \Delta x) &= g(f(x + \Delta x)) = g(f(x) + \Delta y) \\ &= g(f(x)) + g'(f(x))\Delta y + \underbrace{\Delta y \cdot \varepsilon(\Delta y)}_{\textcircled{3}} \\ &= g(f(x)) + g'(f(x))(f'(x)\Delta x + \varphi_1(\Delta x)) \\ &\quad + (f'(x)\Delta x + \varphi_1(\Delta x)) \times \varepsilon(\Delta y) \\ &= g(f(x)) + g'(f(x))f'(x)\Delta x \\ &\quad + \underbrace{g'(f(x))\varphi_1(\Delta x)}_{\textcircled{2}} + \underbrace{(f'(x)\Delta x + \varphi_1(\Delta x))\varepsilon(\Delta y)}_{\textcircled{4}}. \end{aligned}$$

よって, 関数 $\varepsilon(h)$ は

$$\varepsilon(h) = \begin{cases} \frac{\varphi_2(h)}{h} & (h \neq 0) \\ 0 & (h = 0) \end{cases}$$

と選ぶ.

$$\begin{cases} \lim_{h \rightarrow 0} \varepsilon(h) = 0, & \varepsilon(0) = 0. \end{cases}$$

$$g(y+h) = g(y) + g'(y)h + h\varepsilon(h) \quad (|h| < \delta)$$

が成り立つ.